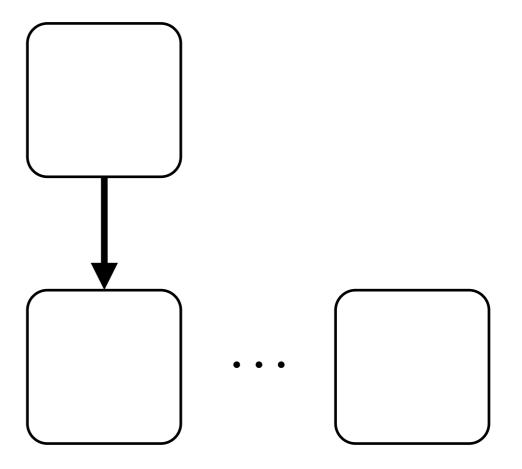
# Guaranteeing memory safety in Rust

Nicholas D. Matsakis Mozilla Research

#### Hashtable in C / C++

```
template<class K, class V>
struct Hashtable {
    Bucket<K,V> *buckets;
    unsigned num_buckets;
template<class K, class V>
struct Bucket {
    bool occupied;
    K key;
    V value;
```



#### Insertion

```
template<class K, class V>
void insert(
 Hashtable<K,V> *table,
  K key,
  V value)
                                 Grow arrays as necessary
  if (table full) {
    table->buckets = realloc(table->buckets, ...);
  unsigned index = find_bucket(table, &key);
  table->buckets[index].key = key;
  table->buckets[index].value = value;
  table->buckets[index].occupied = true;
                                       Initialize the bucket
```

#### Assumptions

```
template<class K, class V>
void insert(
 Hashtable<K,V> *table,
  K key,
  V value)
                                 Buckets array is unaliased
  if (table full) {
    table->buckets = realloc(table->buckets, ...);
  unsigned index = find_bucket(table, &key);
  table->buckets[index].key = key;
  table->buckets[index].value = value;
  table->buckets[index].occupied = true;
                                     Never access bucket if
```

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occupied is false

## Living on the edge

Violating either of these assumptions can lead to a crash

=>

Crashes lead to exploits.

## Privacy?

Q: Doesn't privacy solve this?

A: Not really.

#### Table lookup

```
template<class K, class V>
V* find(
  Hashtable<K,V> *table,
                                     Find the correct bucket.
  K *key)
  unsigned index = find bucket(table, &key);
  if (!table->buckets[index].occupied)
    return NULL;
                                        Bucket initialized?
  return &table->buckets[index].value;
                                     Return pointer into the
```

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bucket array!

#### Vulnerability

# A problem for everyone

Q: Doesn't garbage collection solve this?

A: Not really.

e.g., Java's ConcurrentModificationException

## Iteration holds pointers

```
(for-each hashtable (lambda (x) ...))
```

if hashtable is mutable, same problem

#### Implications

During iteration, the buckets array is aliased.

+

Insertion can resize the array.

=>

Insertion during iteration can allow a foreign website to take over your computer. (e.g., Bug 810718)

#### Enter: Rust



Think: C++ meets ML/Haskell meets Erlang (meets Cyclone, ML Kit, and many others)

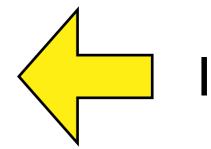
#### Credit where credit is due

Rust is the product of a community too numerous to list here.

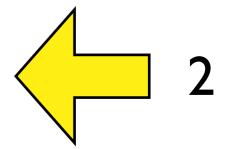
https://github.com/mozilla/rust/blob/master/AUTHORS.txt

#### What Rust has...

- Traits (type classes)
- Ownership (affine types)



- Algebraic data types
- Lifetimes (regions)



Actor-style concurrency

#### What Rust doesn't have...

- Null pointers
- Dangling pointers
- Segmentation faults
- Data races
- Mandatory GC

## Why optional GC?

- I. Efficiency and predictability
- Avoid unpredictable latency
- Particularly on mobile
- 2. Memory disjoint tasks
- If you can send memory, you can free it.

## Preview: Freezing

```
let mut table: Hashtable
Create hashtable
let value = table.find(...);

table.insert(...);

match value { ... }

Illegal, hashtable frozen
table.insert(...);
Illegal, hashtable frozen
```

OK, value is out of scope

#### Hashtable in Rust

```
struct Hashtable<K,V> {
   buckets: ~[Option<Bucket<K,V>>]
}

struct Bucket<K,V> {
   key: K,
   value: V
}

c[T]: Owned pointer
   to an array
```

## What's in a type

~[T]: Owned pointer to an array

~[Option<Bucket<K,V>>]

Option<T>: I guess you've seen this before

#### What is ownership?

In Rust, ownership means:

The Right To Free Memory



Control over aliasing

## Ownership is implicit in C/C++

Some pointers are temporary:

```
memcpy(void *dest, const void *src, uintptr_t count);
V* find(Hashtable<K,V> *table, K *key);
```

#### Some pointers are not:

```
struct Hashtable {
    Bucket<K,V> *buckets;
}

void *realloc(void *, ...);
void free(void *);
```

#### Ownership is explicit in Rust

Temporary pointers are designated &T

Owned pointers are designated ~T

Managed pointers are designated @T

## Owned pointers

Owned pointers never alias one another and are automatically freed:

#### Owned pointers are moved

Owned pointers are moved from place to place.

```
fn move_from() {
  let x: \simint = \sim22;
                  Moves x into callee's stack frame
  move to(x);
  printf("%d", *x); // Error: x was moved
                              x no longer accessible
fn move_to(y: ~int) {
              Freed by callee
```

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#### Moving into a data structure

Owned pointers can also be owned by a data structure.

Only accessible via table

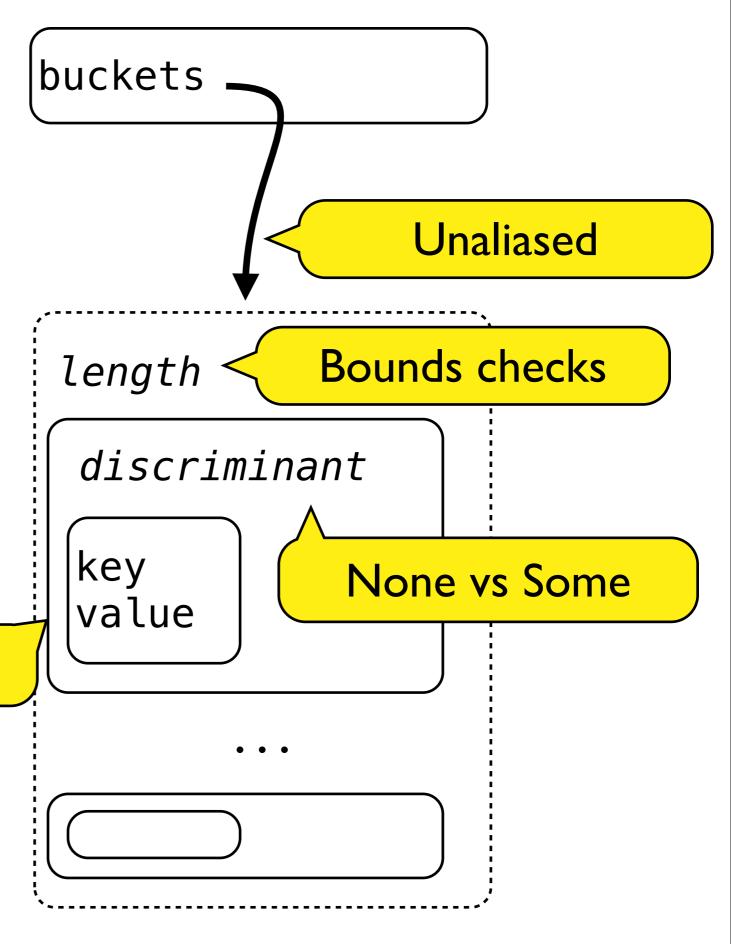
#### Hashtable<K,V>

~[Option<Bucket<K,V>>]

Option<Bucket<K,V>>

Bucket<K,V>

Data if Some, like a C union



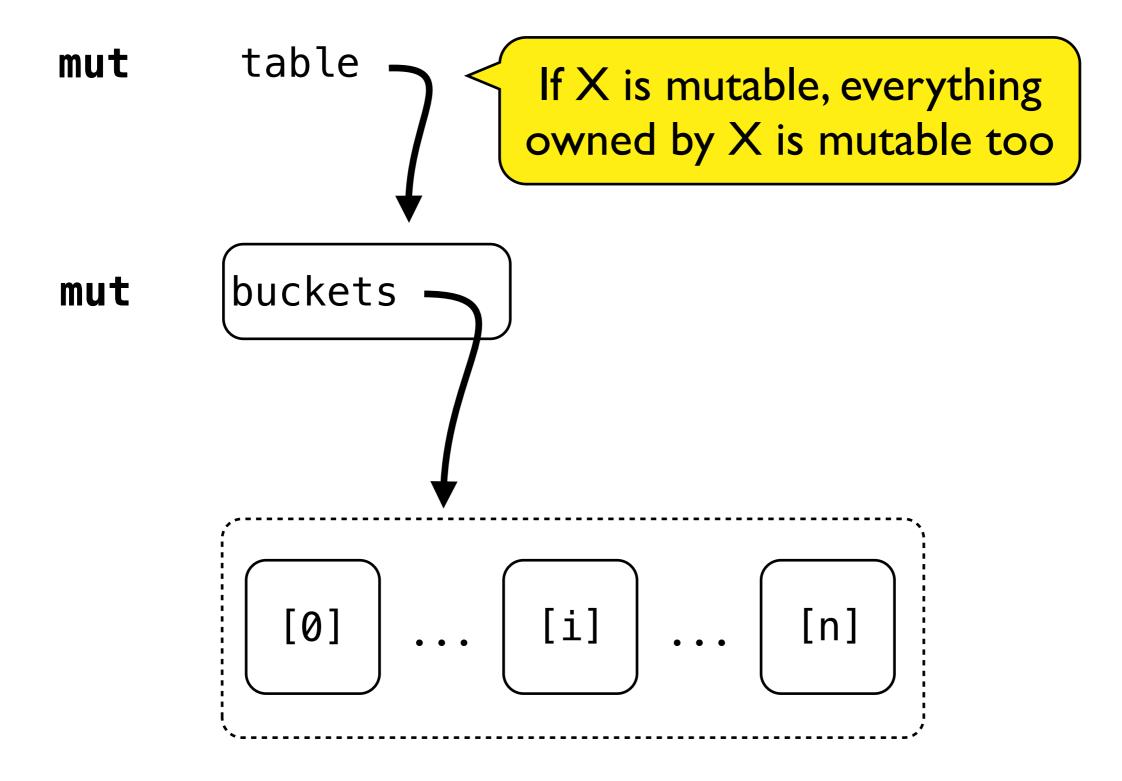
#### Insertion, first attempt

```
fn insert<K,V>(
  mut table: ~Hashtable<K,V>,
  key: K,
                Take ownership of a hashtable
  value: V)
  -> ~Hashtable<K,V>
  if (table full) {
    table.buckets = resize_buckets(table.buckets);
                               Safe because of ownership
  let index = find_bucket(table, &key);
  table buckets[index] = Some(Bucket {key: key,
                                         value: value});
                    Update buckets array
  return table;
               Give table back to the caller.
```

## Mutability in Rust

```
let x = 22;
x += 1; // Error
                    Local variables must be
let mut x = 22;
                     declared as mutable
x += 1; // 0K
fn update(x: int) {
  x += 1; // Error
fn update(mut x: int) {
  x += 1; // 0K
                      Parameters too
```

## Mutability and ownership



#### Inherited mutability

```
fn insert<K,V>(
  mut table: ~Hashtable<K,V>,
  key: K,
                       table is mutable
  value: V)
  -> ~Hashtable<K,V>
  if (table full) {
    table.buckets = resize_buckets(table.buckets);
                                table.buckets is mutable
  let index = find_bucket(table, &key);
  table.buckets[index] = Some(Bucket {key: key,
                                         value: value});
           table.buckets[index] is mutable
  table
```

#### Ownership is explicit in Rust

Temporary pointers are designated &T

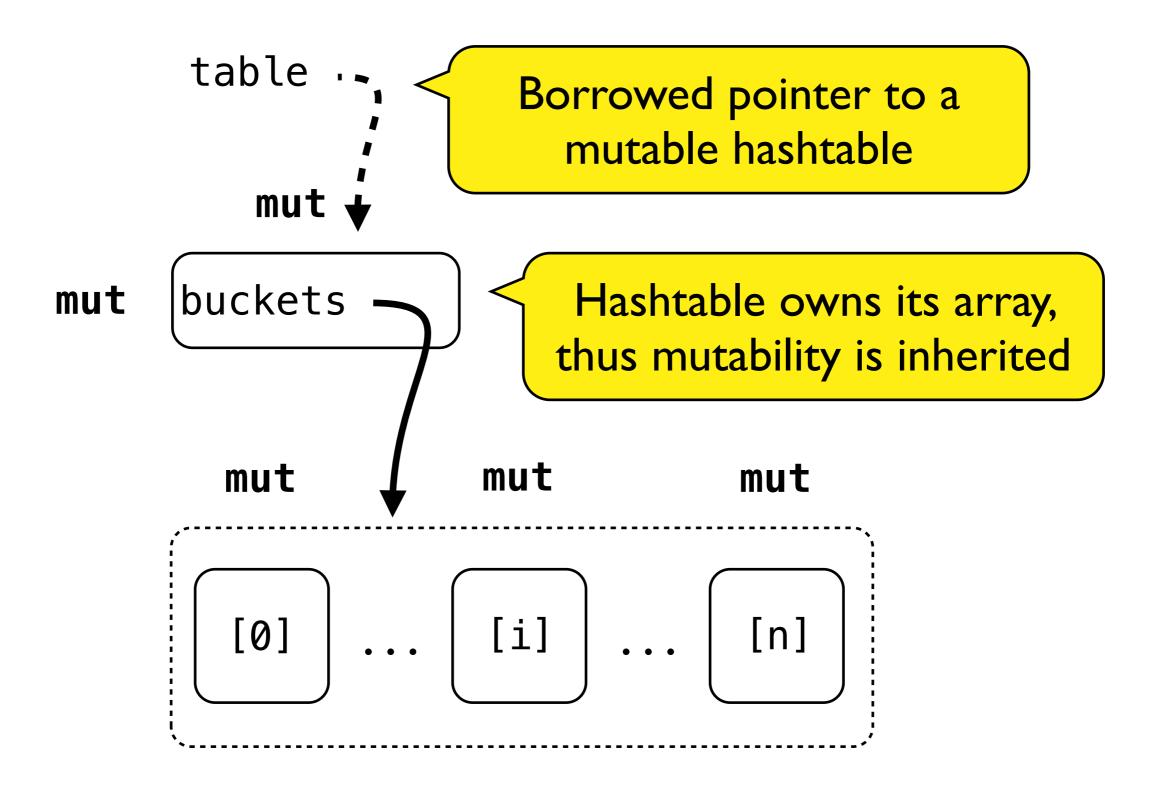
Owned pointers are designated ~T

Managed pointers are designated @T

#### Borrowing

```
fn insert<K,V>(
  table: &mut Hashtable<K,V>,
  key: K,
                Mutable borrowed pointer as input
  value: V)
                              Mutable borrow
  if (table full) {
    resize buckets(&mut table.buckets);
                                      Immutable borrows
 let index = find_bucket(&*table, &key);
 table.buckets[index] = Some(Bucket {key: key,
                                        value: value});
```

## Borrowing



# Mutability and uniqueness

&mut T Borrowed pointer to mutable data

Type system guarantees uniqueness: no other mutable pointer to the same data.

Linearly tracked.

#### Mutable borrows

```
Mutable local variable
let mut x = 5;
                    x is borrowed for the lifetime of y
  let y = \& mut x;
  x += 1; // Error: borrowed
  *y += 1; // 0K
                 Borrow has expired, y out of scope
x += 1; // 0K
                          Mutable borrows prevent
                            owner from writing.
```

#### Lifetimes

```
let mut x = 5;
{
    let y = &mut x;
    x += 1; // Error: borrowed
    *y += 1; // OK
}
x += 1; // OK
```

### Lifetime parameters

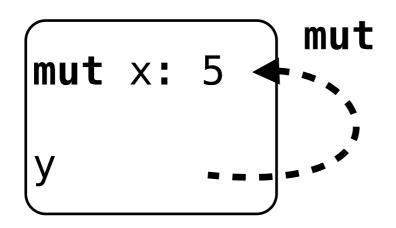
```
fn insert<K,V>(
  table: &mut Hashtable<K,V>,
  key: K,
  value: V)
             For any lifetime 'a specified by caller...
fn insert<'a,K,V>(
  table: &'a mut Hashtable<K,V>,
  key: K,
  value: V)
```

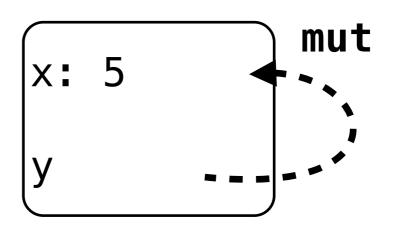
#### I. Data must be mutable

```
let mut x = 5;
let y = &mut x;
let y = &mut x;

OK

ILLEGAL
```





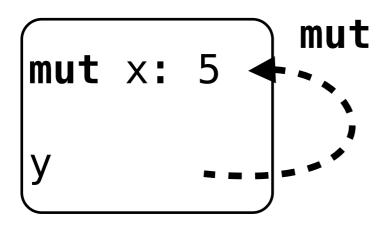
2. Data must live long enough

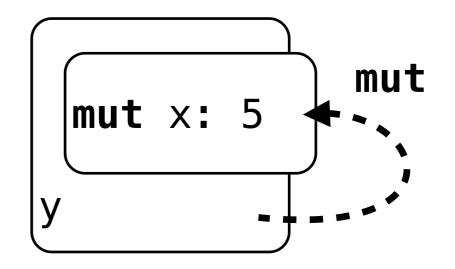
```
{
   let mut x = 5;
   let y = &mut x;
}
```

let y;
{
 let mut x = 5;
 y = &mut x;
}
\*y += 1;

OK

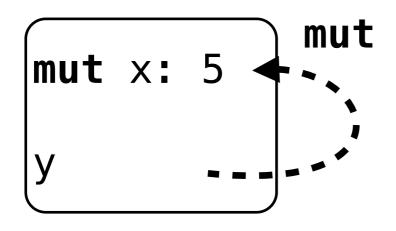


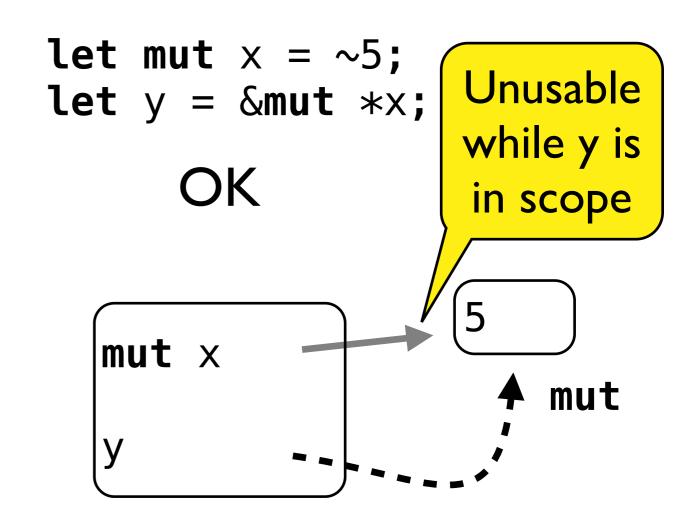




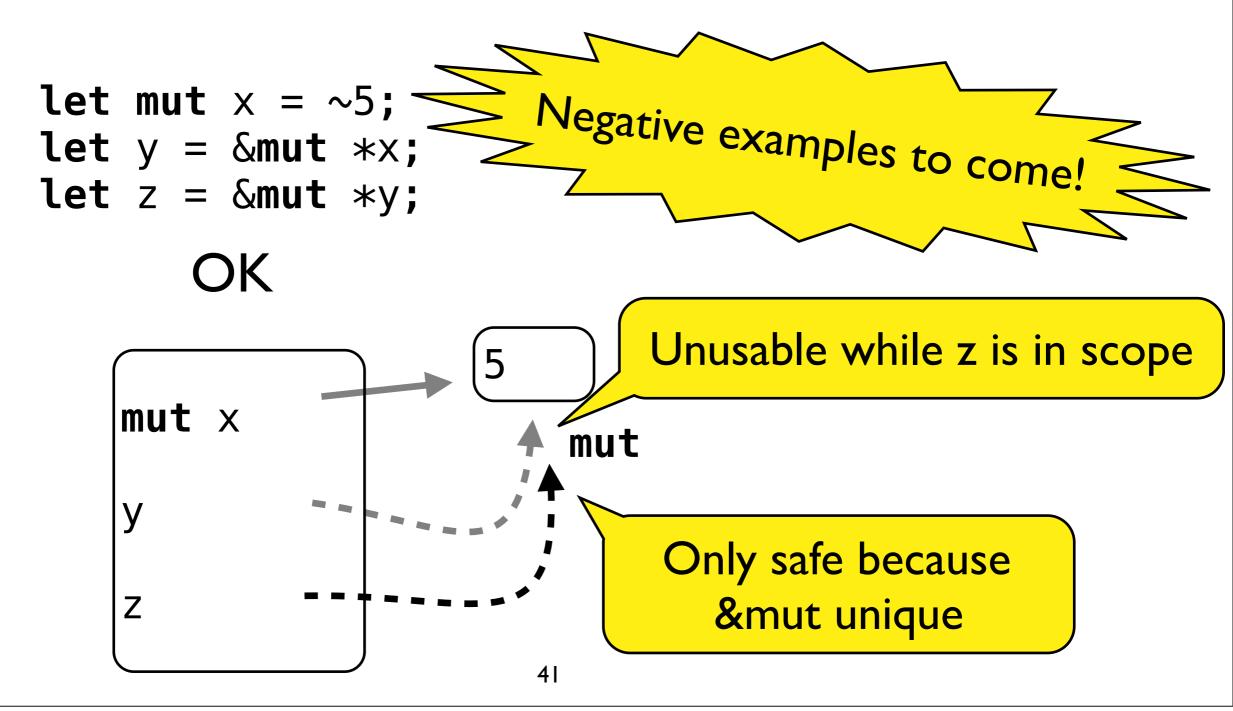
3. Data must be uniquely referenced

```
let mut x = 5;
let y = &mut x;
OK
```





3. Data must be uniquely referenced (cont'd)



# Immutable borrowed pointers

**&T** 

Borrowed pointer to immutable data

May be aliased

Stronger than C++ const, **nobody** can modify it

#### Immutable borrows

```
let x = 5;
{
  let y = &x;     Original can still be read
  print(x);
}
```

Immutable borrows do not restrict owner from reading.

But there are other restrictions...

# Freezing

```
tet mut x = 5;
{
    let y = &x;
    x += 1; // Error
}
x declared as mutable

Frozen for lifetime of y

y out of scope, unfrozen

y out of scope, unfrozen
```

I. Data must live long enough

```
{
  let x = 5;
  let y = &x;
}
```

```
let y;
{
   let x = 5;
   y = &x;
}
print(*y);
```

OK

**ILLEGAL** 

Same as with &mut

2. Mutable data must be uniquely referenced.

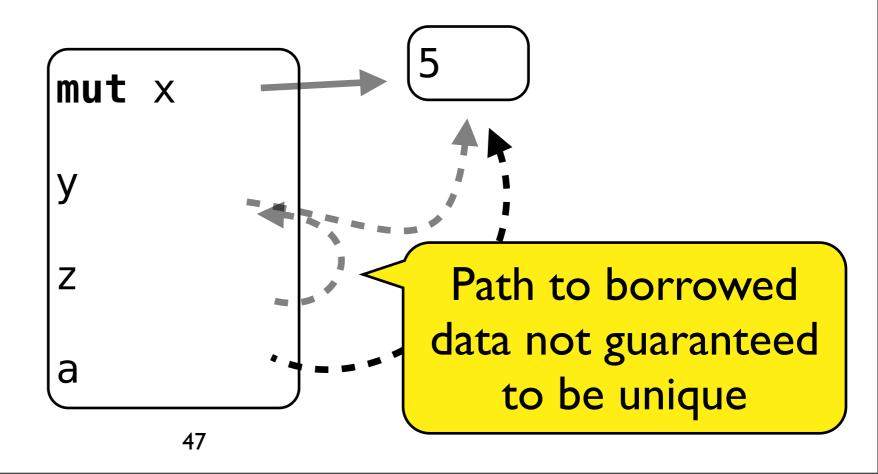
```
let mut x = 5;
let y = &x;
let y = &*x;
OK
```

```
let mut x = \sim 5;
let y = \& mut *x;
let z = \& *y;
```

Same as with &mut

# An illegal case...

```
let mut x = ~5;
let y = &mut *x;
let z : & &mut int = &y;
let a = &mut **z;
```



# An illegal case (cont'd)

```
let mut x = \sim 5;
let y = \&mut *x;
let z : \& \&mut int = \&y;
let z2 = z;
let a = \&mut **z;
       mut x
       z2
                           *a and **z2 access
                            the same data!
       a
              48
```

# Let's put it all together

Remember find()?

Returns a pointer into the hashtable.

Goal: Ensure that hashtable is not modified while this pointer is live.

# Lifetime parameters in find

```
fn find<'a,K,V>(
   table: &'a Hashtable<K,V>,
   key: &K)
   -> Option<&'a V>
{
   ...
}
```

# Lifetime parameters in find

```
Given: pointer with lifetime 'a to an immutable hashtable

fn find<'a,K,V>(
    table: &'a Hashtable<K,V>,
    key: &K)
    -> Option<&'a V>

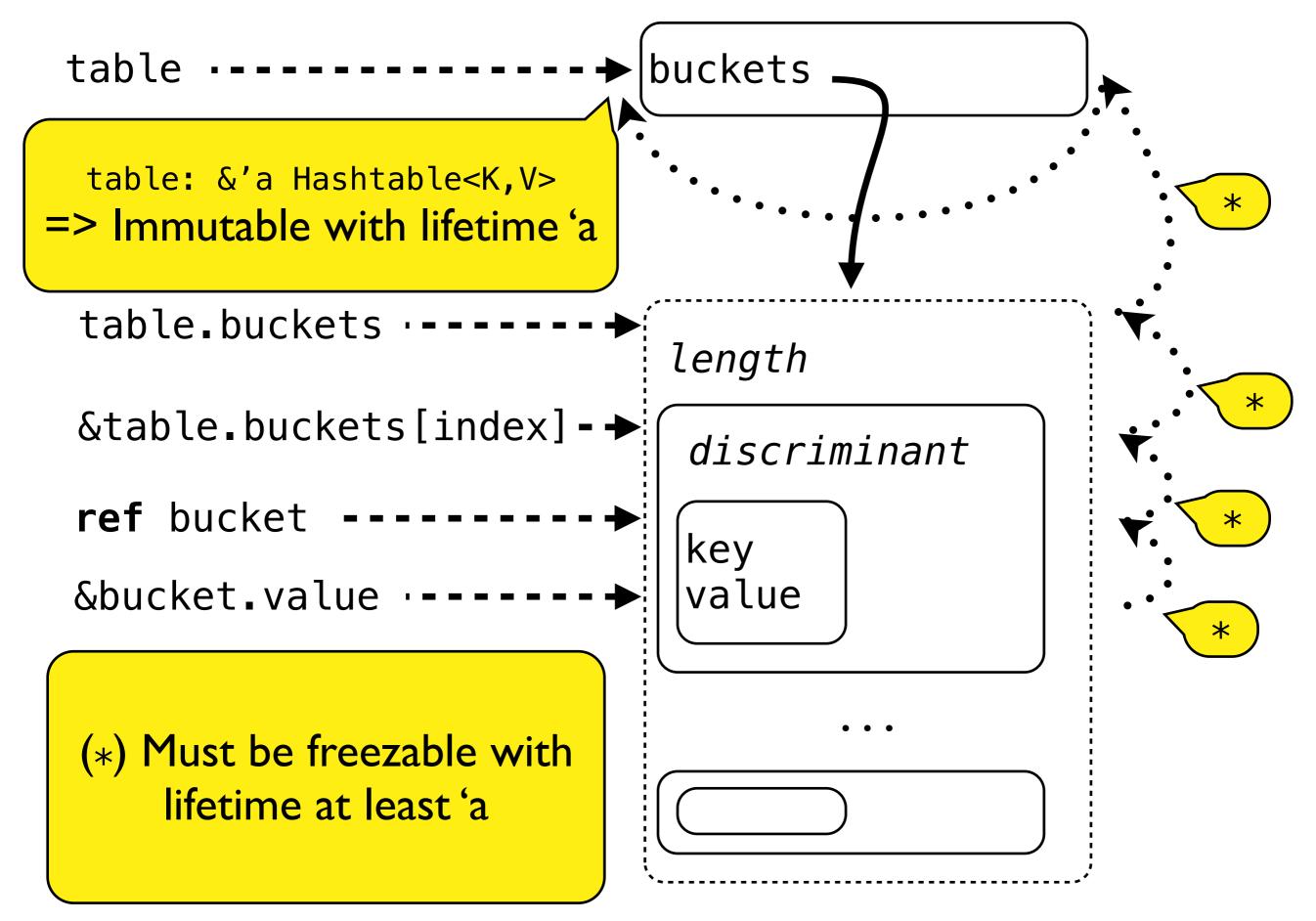
{
    Yields: (optional) pointer with lifetime 'a to an immutable value V
```

In other words, the value returned is valid as long as:

- I. the hashtable is valid
- 2. the hashtable is not mutated

#### Find

```
fn find<'a,K,V>(
  table: &'a Hashtable<K,V>,
  key: &K)
                      Search for index...
  -> Option<&'a V>
  let index = find_bucket(table, &key);
  match table.buckets[index] {
    None => None,
    Some(ref bucket) => Some(&bucket.value)
           Creates a pointer into value being matched
```



#### Voila

```
let mut table: Hashtable<K,V> = ...;
  let value = find(&table, ...);
  insert(&mut table, ...); Frozen for lifetime of value
  match value { ... }
                                Error, frozen.
insert(&mut table, ...);
              OK, value out of scope.
```

#### But wait, there's more...

```
fn compute(input: &[T], output: &mut [T]) {
  if output.len() >= threshold {
    let (left, right) = output.split();
    parallel::do([
          compute(input, left),
          compute(input, right)]);
  } else {
          Immutable, safe to share.
                Mutable but disjoint. Safe.
                    Divide buffer into two disjoint halves
```

# Recap #1

Owned pointers are moved, not freely aliased.

 $\Rightarrow$ 

Safe to free.

Safe to send to another thread.

Safe to resize.

### Recap #2

Borrowing limits the owner for the lifetime of the borrow.

 $\Rightarrow$ 

Borrowed values cannot be moved or sent between tasks.

Mutable borrowed values can be temporarily frozen.

### Recap #3

Mutable borrowed pointers are unique.



Allows reborrowing, resizing, and possibly parallelism beyond actors (wip).

#### More information?

Nicholas Matsakis nmatsakis@mozilla.com

rust-lang.org

smallcultfollowing.com/babysteps